

On Stable Marriages and Greedy Matchings

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Background

The Stable Marriage (SM) problem has a long and rigorous history.

Greedy Matchings (GM) have applications in CSC applications.



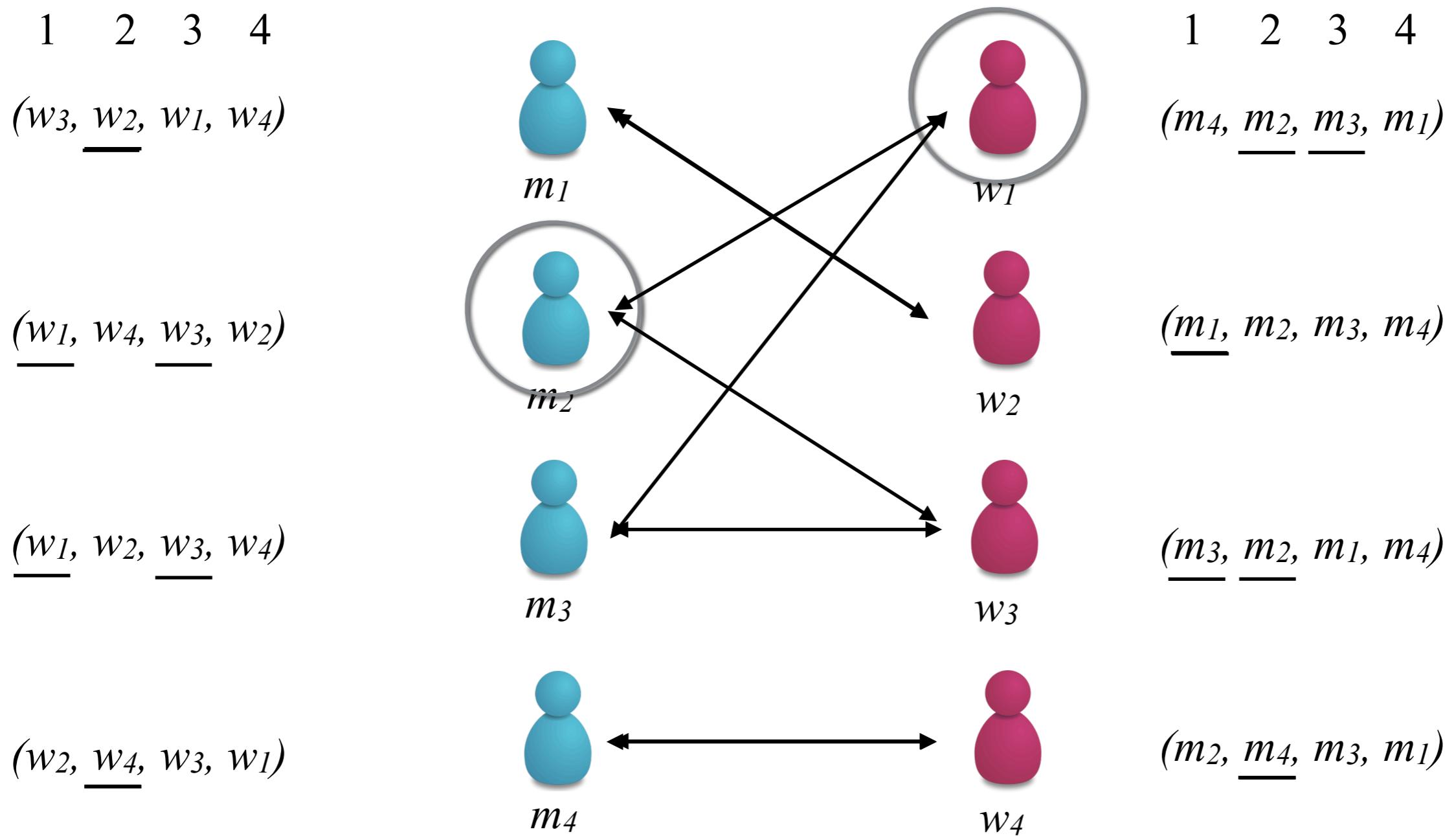
Objective:

- ▶ Formalize the connections between the Stable Marriage problem and computing Greedy Matchings.

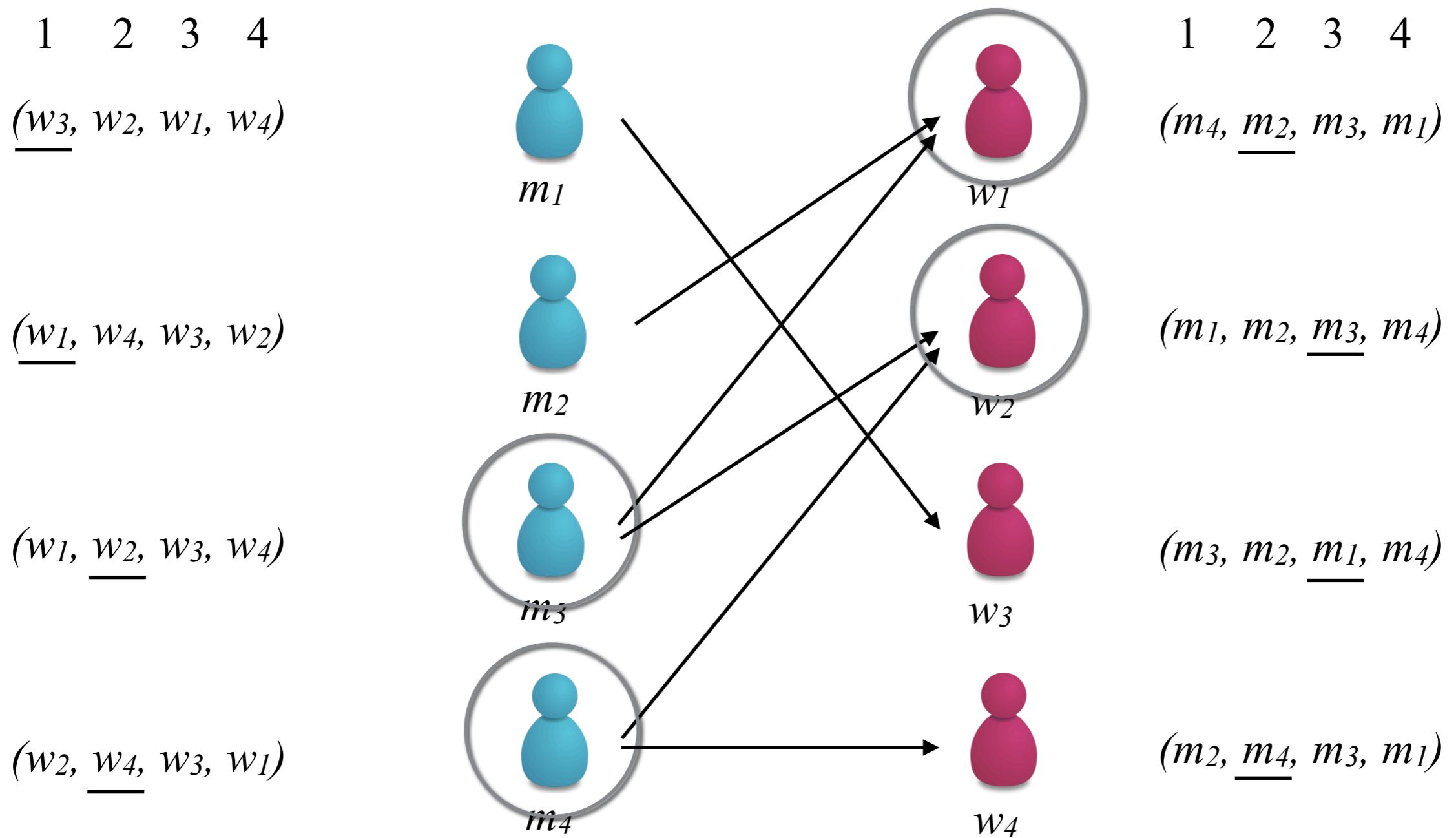
Consequences:

- ▶ Many "new" algorithms for computing GM are variants of algorithms for SM.
- ▶ Parallel algorithms for computing GM can also be applied to SM.

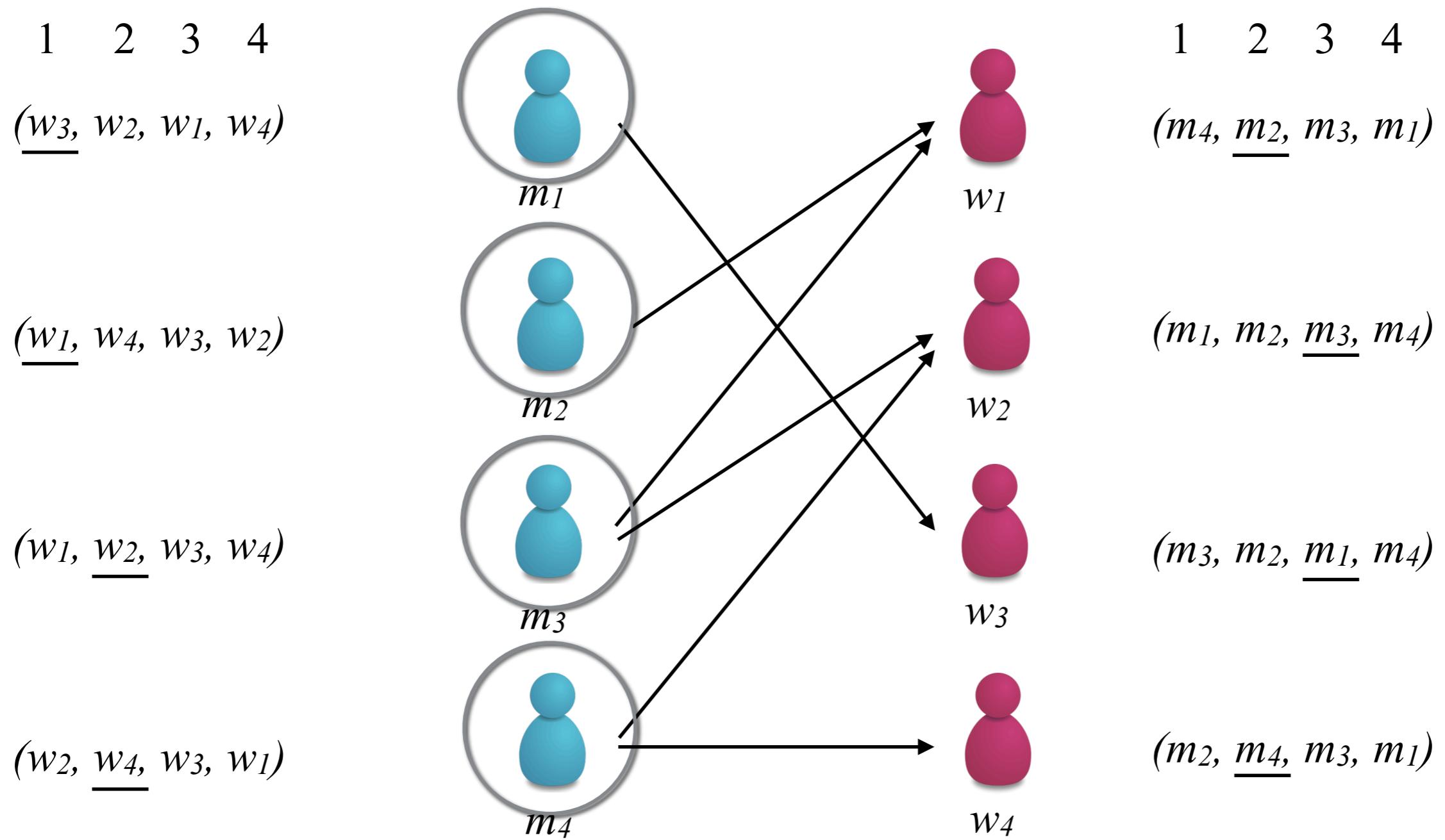
The Stable Marriage Problem



The Gale-Shapley Algorithm

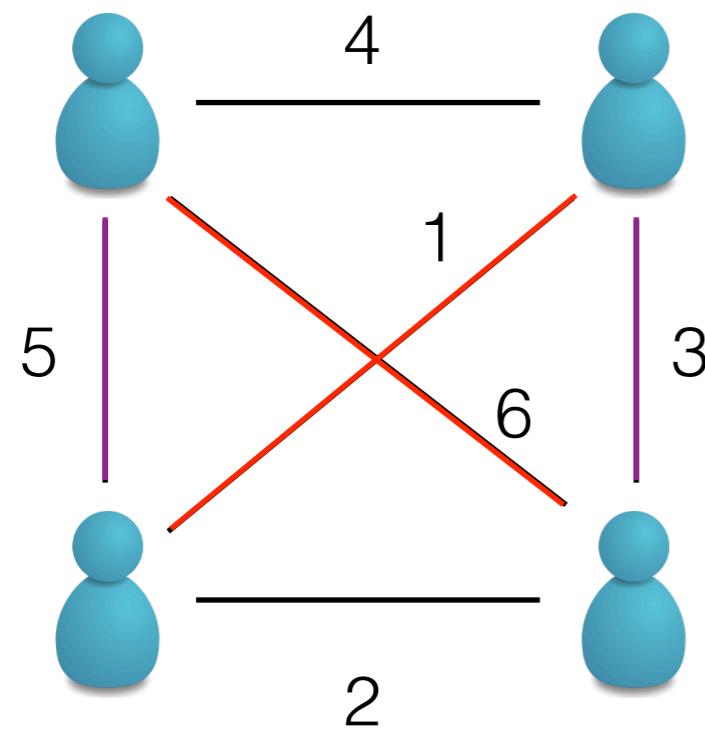


The McVitie-Wilson Algorithm



Implementation: GS uses a queue while MW uses a stack for the remaining men

Greedy Matching



$$M = \emptyset$$

While there are edges remaining

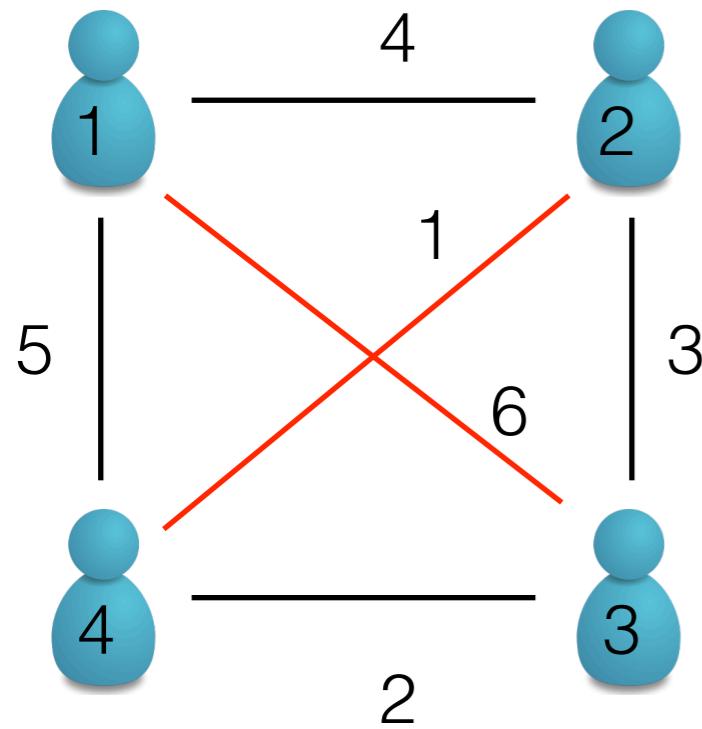
e = heaviest edge

$$M = M \cup \{e\}$$

remove edges incident on e

$$w(\text{greedy}) \geq 0.5 w(\text{optimal})$$

Computing a Greedy Matching using a Stable Marriage algorithm

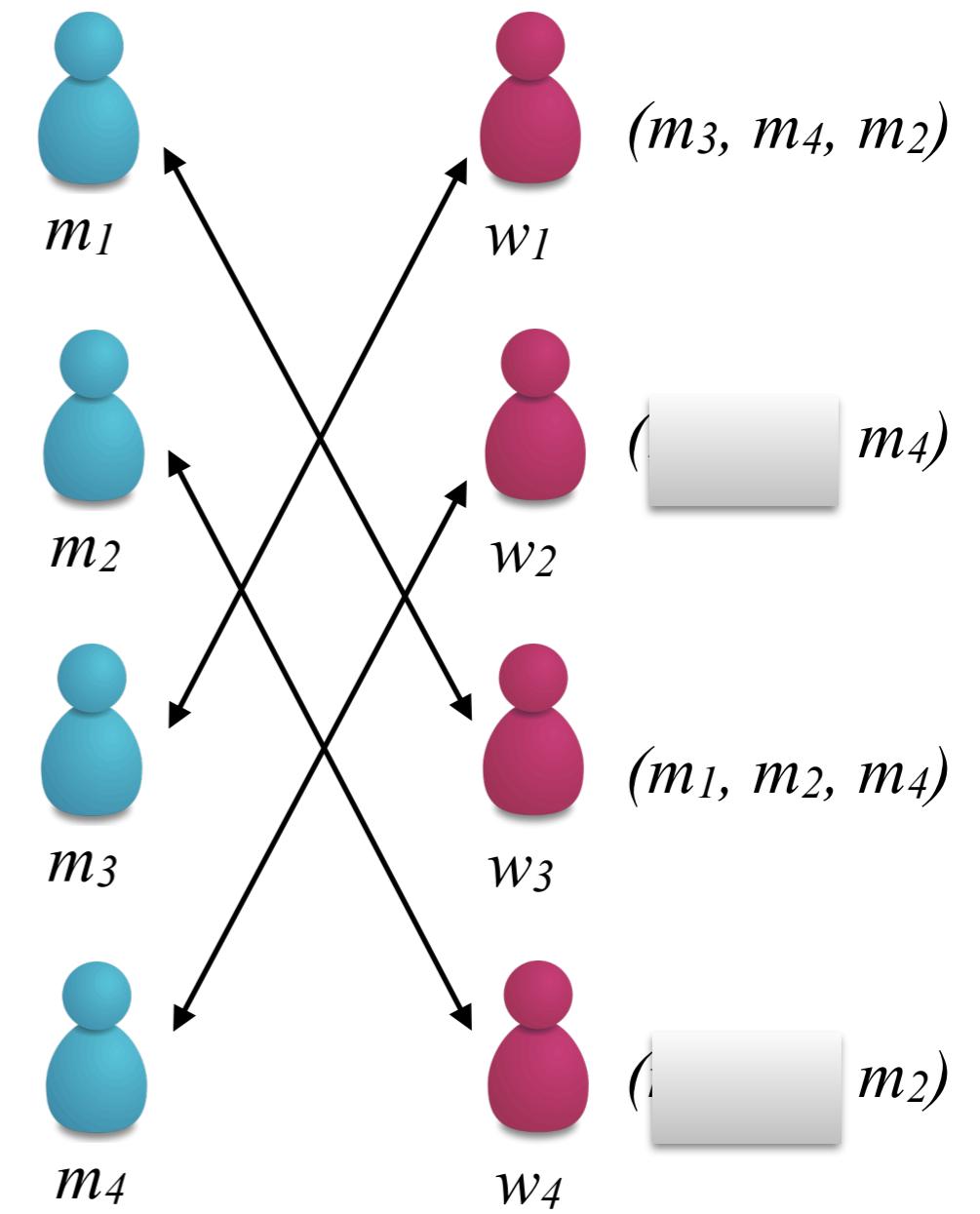


(w_3, w_4, w_2)

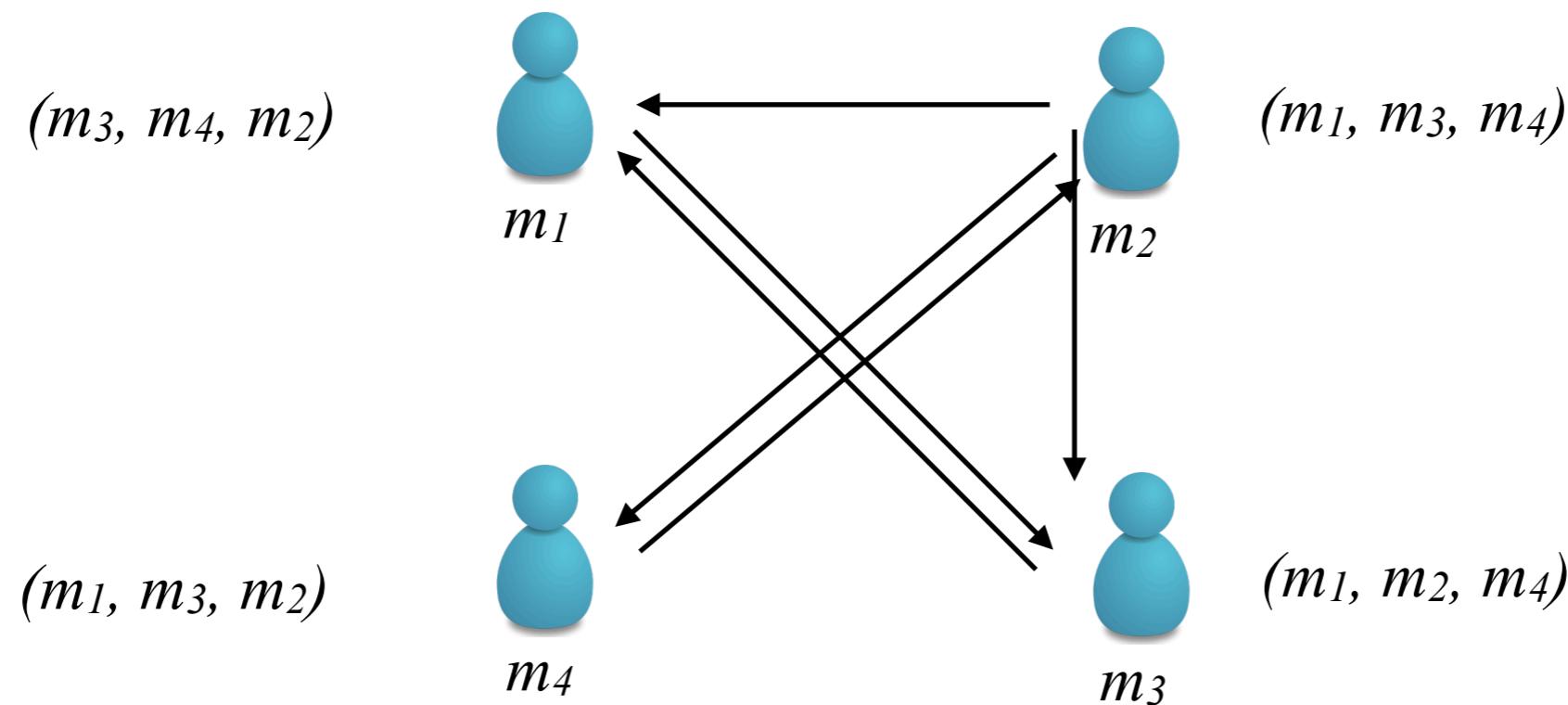
(\quad, w_4)

(w_1, w_2, w_4)

(\quad, w_2)



Using McVitie-Wilson directly

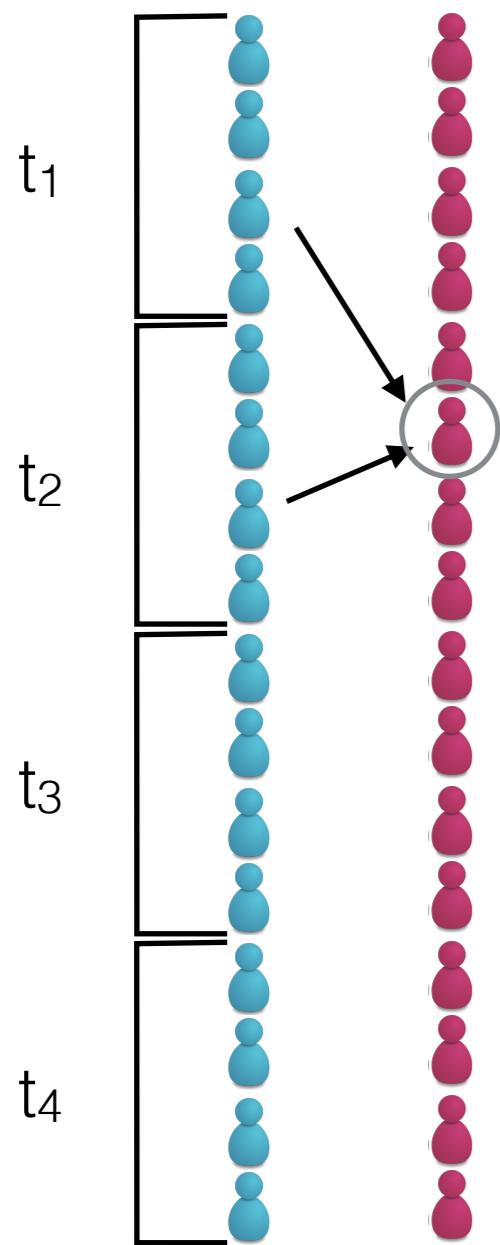


This is exactly the suitor-algorithm [Manne & Halappanavar 14]

which builds on:

- ▶ The pointer algorithm [Manne & Bisseling 08]
- ▶ The Preis algorithm [Preis 99]

Going parallel



Threads run either the Gayle-Shapley or the McVitie-Wilson algorithm

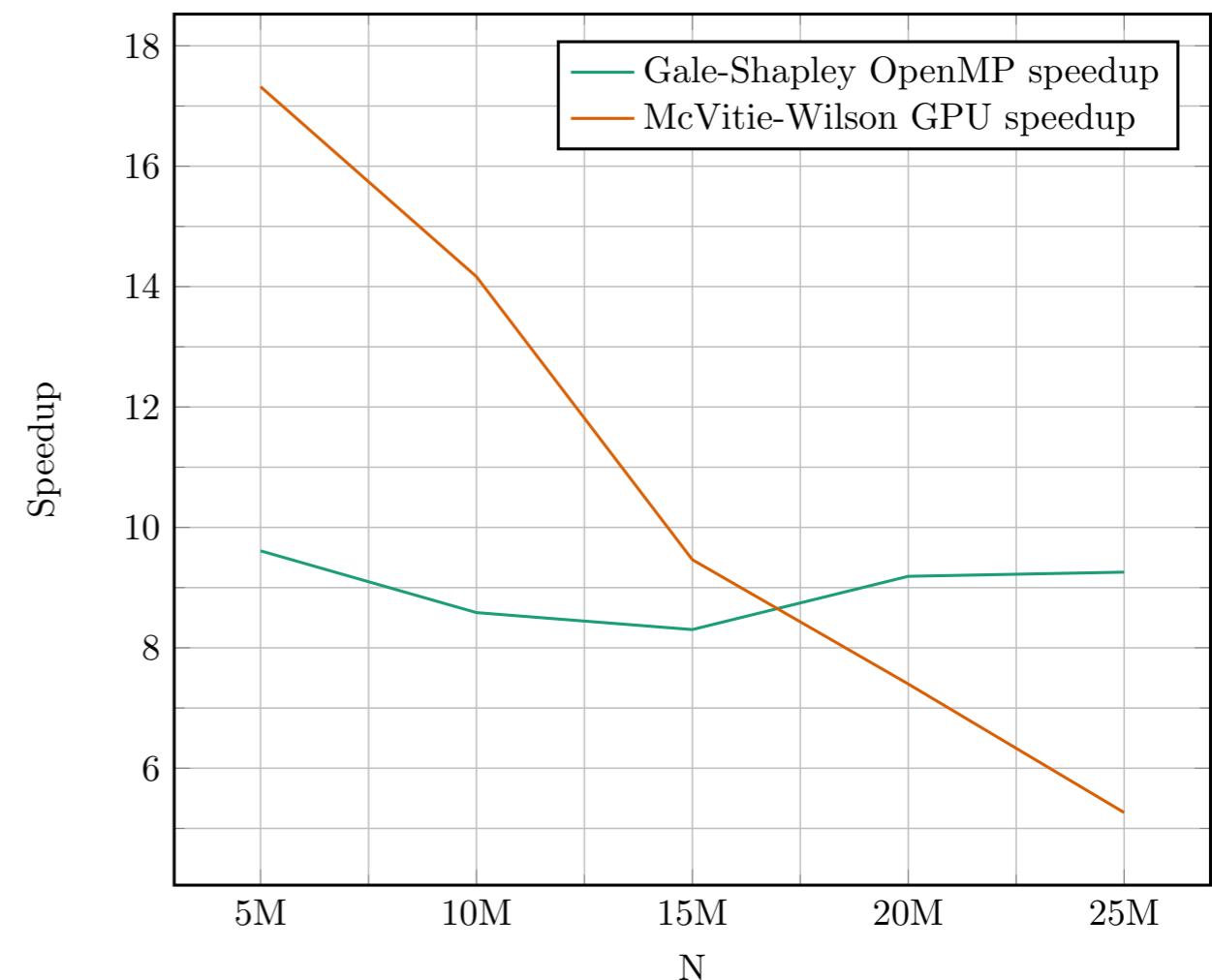
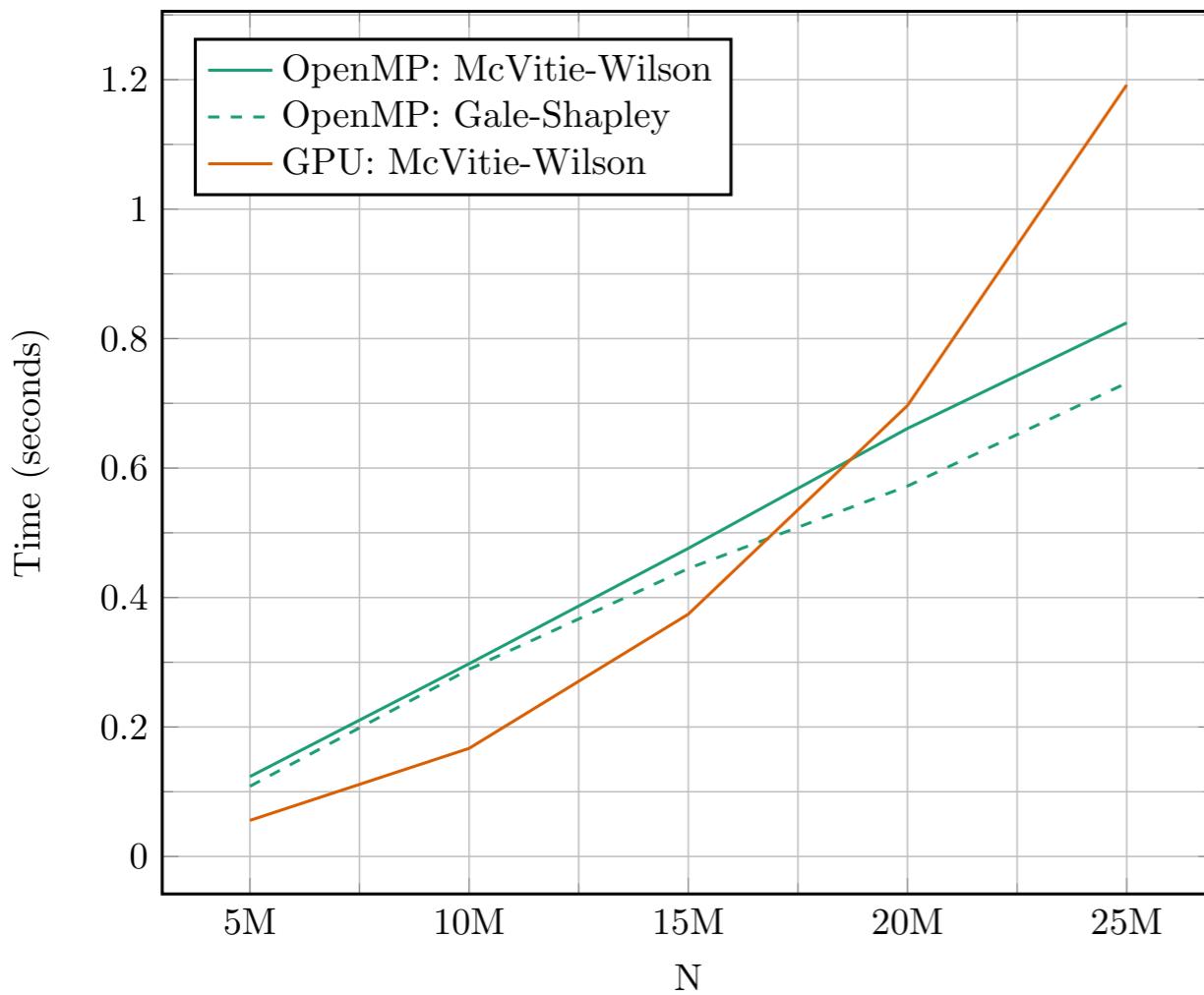
Use compare-and-swap to protect "women"

Implementations using both OpenMP and GPU

“Easy” problems

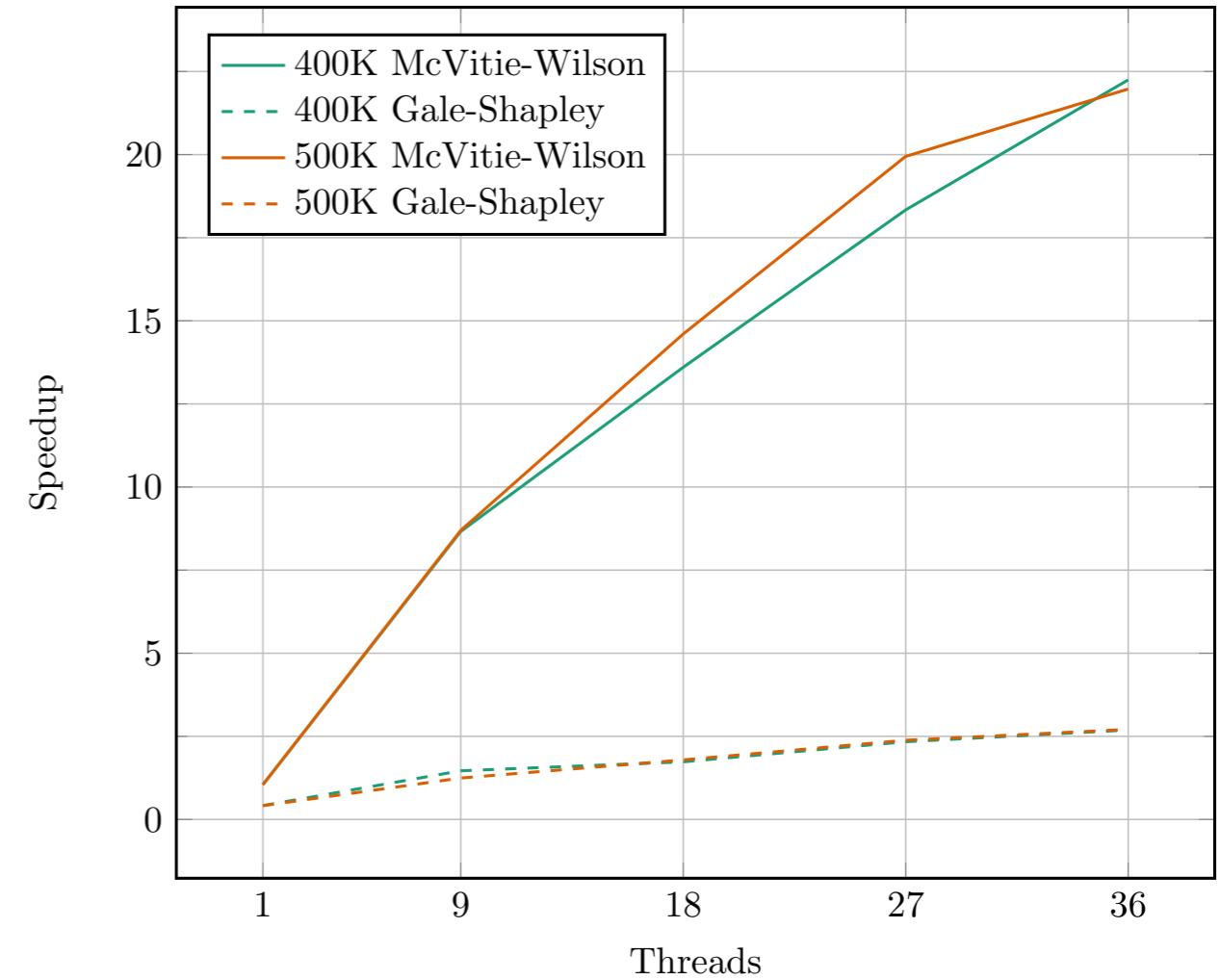
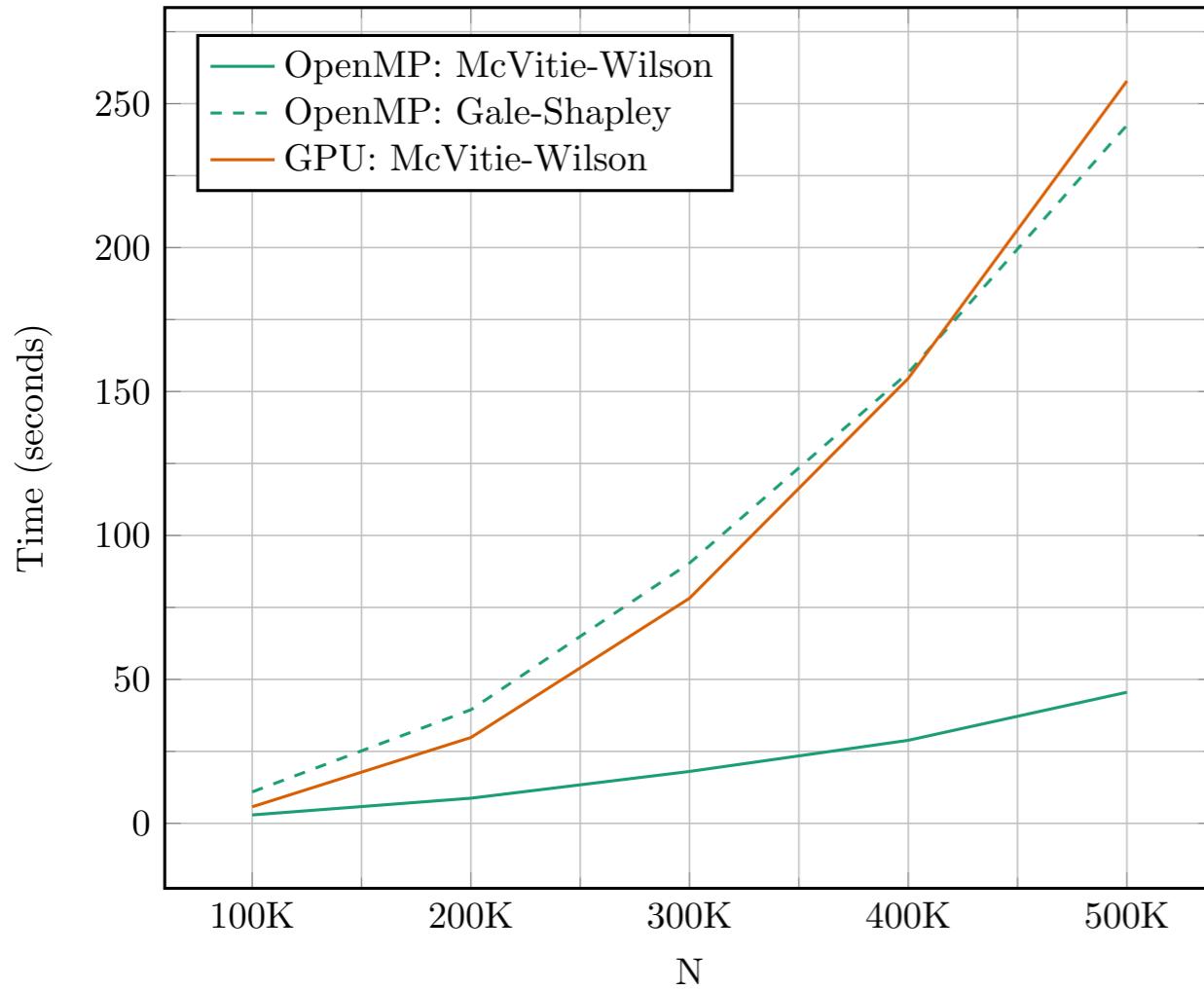
- Each man selects between $\log n$ and $2 \log n$ women and ranks randomly.
- Women rank men who rank them (randomly).
- Expected total work $n \log n$

OpenMP: Using 36 threads on two Intel Xeon E5-2699 processors
GPU: Tesla K40m with 2880 cores



“Hard” problems

- Each man uses the same total ranking of all the women
- Each woman uses the same total ranking of all the men.
- Expect high contention for the same “women”
- Total work will be $(n+1)n/2$



Concluding Remarks

- Recent greedy b-matching algorithms also follows directly from algorithms for the many-to-many stable assignment problem.
- Major open question [Manlove 13]: Is Stable Marriage in NC?
 - Maybe not so relevant...
- Stable Marriage assumes sorted priority lists, whereas Greedy Matching makes no such assumption.
 - Preis solved Greedy Matching in $O(m)$ time.
 - Can Stable Matching with unsorted weighted priority lists also be solved in $O(m)$ time?

Gale-Shapley implementation

Place all vertices in queue Q
while $Q \neq \emptyset$

$u = Q.\text{first}()$

$p = \text{nextCandidate}(u)$

 while $r_p(u) > r_p(\text{suitor}(p))$

$p = \text{nextCandidate}(u)$

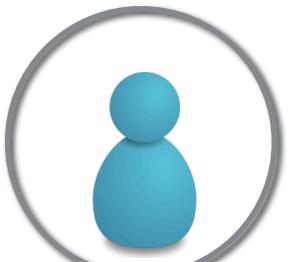
 if $\text{suitor}(p) \neq \text{null}$

$Q.\text{add}(\text{suitor}(p))$

$\text{suitor}(p) = u$

(w_3, w_2, w_1, w_4)

$p \quad p \quad p$



u

(m_4, m_2, m_3, m_1)

p



$\text{suitor}(p)$

